

Siemens Corporate Technology | August 2015

Hard Partitioning for Linux: The Jailhouse Hypervisor

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Agenda

Motivation

Jailhouse introduction & philosophy

Current status (with demo)

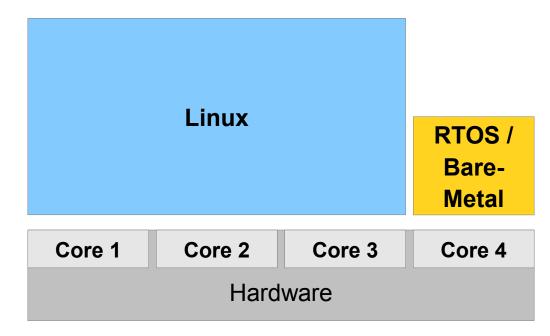
Configuration

Running multiple Linux instances (with demo)

Summary



Asymmetric Multi-Processing (AMP) & Linux





AMP Drivers

Low latency & high throughput

Hard real-time

Preexisting software

Mixed criticality



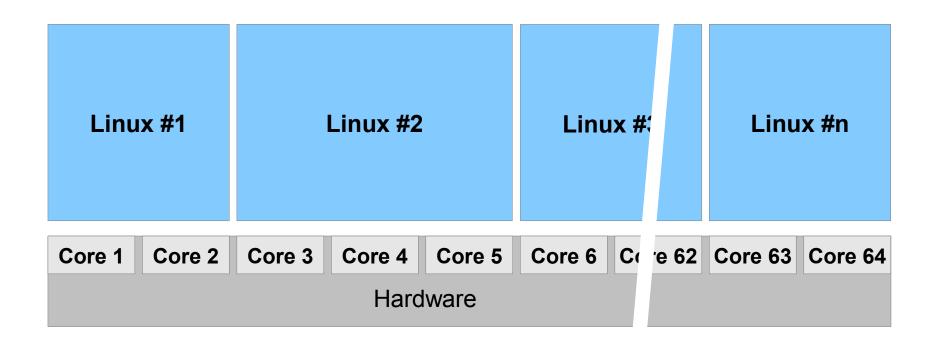








AMP for Linux?





Further AMP Drivers

- Consolidation of (Linux) services
- Enabling of open platforms
- Low latency & high throughput
- Avoid SMP scalability bottlenecks
- Strict quality of service needs



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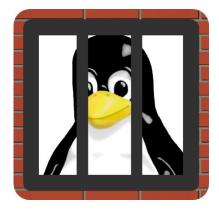
What is Jailhouse?

A tool to run

- ... real-time and/or safety tasks
- ... on multicore platforms (AMP)
- ... aside Linux

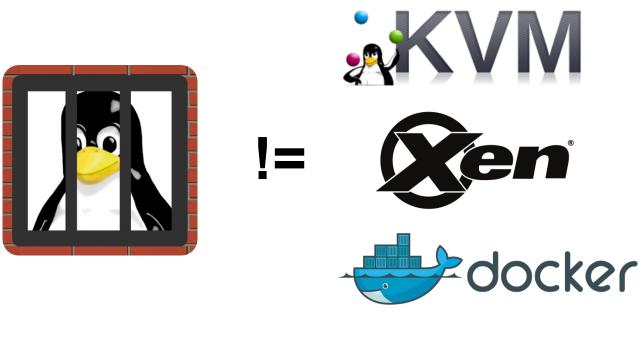
It provides

- strong & clean isolation
- bare-metal-like performance & latencies
- no reason to modify Linux (well, almost)
- ... and it's open source (GPLv2)



What is it not?





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What makes Jailhouse different?

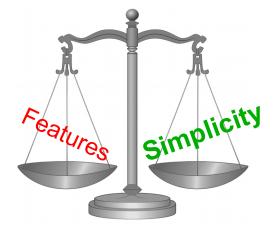
• Use virtualization for isolation – *ok, nothing new*

Prefer simplicity over features

- Resource access control instead of resource virtualization
- 1:1 resource assignment instead of scheduling
- Partition booted system instead of booting Linux
- Do not hide existence of Jailhouse

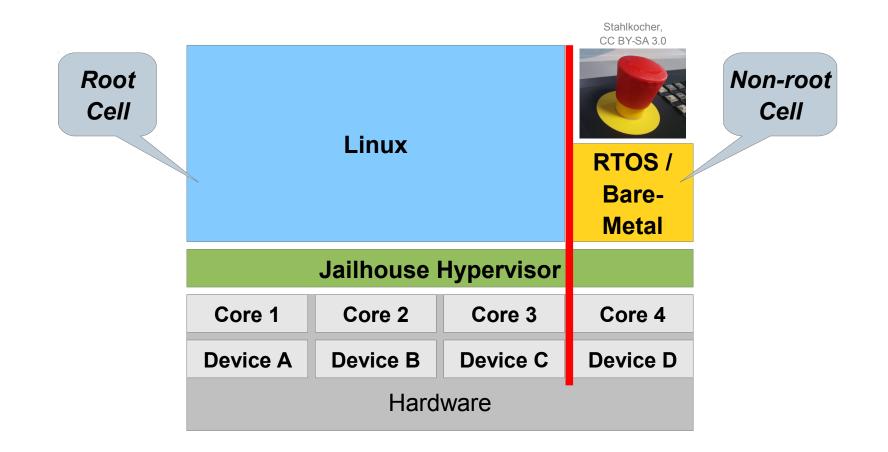
Offload work to Linux

- System boot
- Jailhouse and partition ("cell") loading & starting
- Control and monitoring





AMP with Jailhouse





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Jailhouse Status – x86

- Initial focus on x86, first Intel, then AMD
 - Requirement: VT-x / VT-d, AMD-V
 - AMD IOMMU support to be published soon
- It's small!
 - Currently ~8.5K lines of code (for Intel)
- Direct interrupt delivery
 - Zero VM exits, minimal latencies feasible
 - Max. timer IRQ latency (Xeon D-1540):
- Cache Allocation Technology
 - Intel feature for partitioning L3 cache
 - Code ready to be merged, measurements pending



<2.5 µs



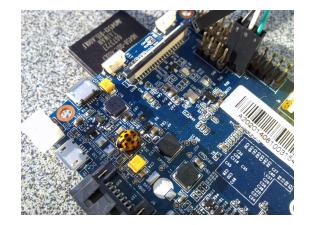
Jailhouse Status – ARM

• ARMv7

- Runs in FastModel, on Banana-Pi, NVIDIA Jetson TK1
- SMMU on to-do list
- It's small too!
 - Currently ~6.5k lines of code (for TK1)

• ARMv8

- First patches just posted by Huawei
- Not yet working, but progressing quickly
- Target: ARMv8 Foundation Model





Live Demonstration, Part I

Running Jailhouse in a virtual machine?!

			Timed Event Loop		
Jailhouse Hypervisor					
Core 0	Core 1	Core 2	Core 3		
Devices			PM UART Timer 2		
QEMU/KVM					



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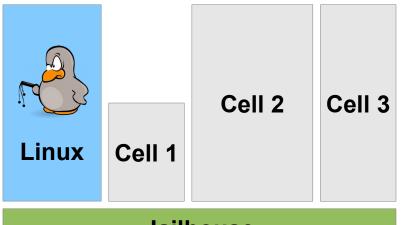
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Two Management Models

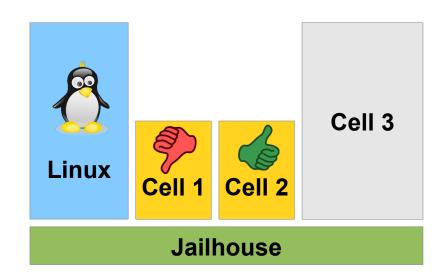
Open Model



Jailhouse

- Linux (root cell) is in control
- Cells not involved in management decisions
- Sufficient if root cell is trusted

Safety Model



- Linux controls, but...
- Certain cells are configured to vote over management decisions
- Building block for safe operation

Configuration

Requirements on raw format

- Easy processing by hypervisor
- Fine-grained control

System configuration => for starting Jailhouse

- Hypervisor reservation
- Relevant platform resources (ACPI / device tree on diet)
- Root cell description

Cell description => for creating cells

Available resources (CPUs, memory, PCI devices, IRQ lines, cache, …)

Creation process

- jailhouse config create my-system.c
- Manual review / post-processing
- Compile: my-system.c \rightarrow my-system.cell
- Cell config: derive from system config



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Future of Jailhouse Configuration

In a nutshell

- Precise & flexible
- ...but not yet convenient

Improvements under discussion

- Cell config generation
- Better consistency checks
- Format revisiting: rebase over device tree?
- ARM will drive this, but x86 may join
- Based on improved low-level automation
 - libvirt support

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Linux As Non-Root Cell?

"Can you also run Linux partitions?"

"Would be too much patching of Linux."

"Why not using Xen PV interfaces? Linux already supports it."

"I try to get a linux kernel run as inmate on Jailhouse [on ARM]."

OK, let's think about this again...

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Looking at x86...

Regular PC

- All devices available (APIC, IOAPIC, PIT, PIC, UARTs, CMOS, RTC, PCI bridges...)
- Resource discovery via ACPI
- Sophisticated handover from BIOS / UEFI

Jailhouse non-root cell

- Only subset of devices
 - APIC (with certain addressing restrictions) → IPI, MSI, timer
 - IOAPIC pins (if assigned)
 - UART (if assigned)
 - PM Timer (as reference clock)
 - TSC (as fast-past clock)
 - Selected PCI devices, no bridges
- Basic resource discovery via communication region
 - PM timer address
 - Number of logical CPUs



A Simpler Solution Than Expected

No!

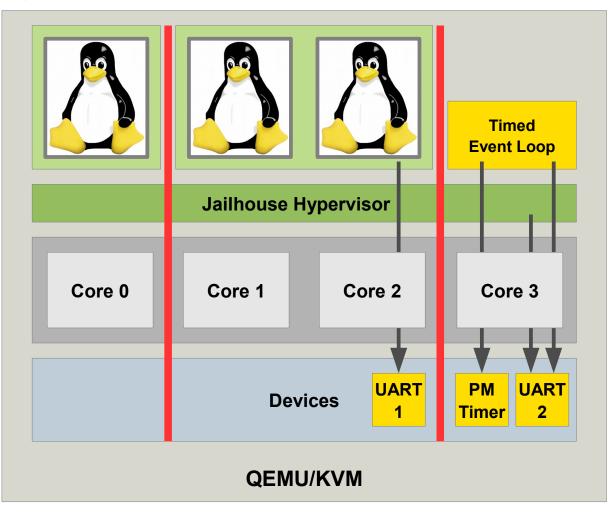
So, we *do* need to patch a lot?

arch/x86/Kconfig	10 +			
arch/x86/include/asm/hypervisor.h	1 +			
arch/x86/include/asm/jailhouse_para.h	27 ++			
arch/x86/kernel/Makefile	2 +			
arch/x86/kernel/apic/apic_flat_64.c	12 +-			
arch/x86/kernel/cou/hypervisor.c	3 +			
arch/x86/kernel/jailhouse.c	229 +++++++++++			
arch/x86/kernel/smpboot.c 7 +-				
<pre>8 files changed, 286 insertions(+), 5 deletions(-)</pre>				



Live Demonstration, Part II

Jailhouse booting Linux

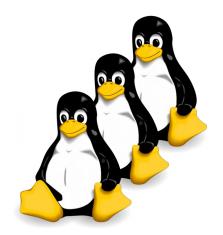




Status of Non-Root Linux Cells

• x86

- Assignment of MSI/MSI-X PCI devices working, no legacy INTx
- SMP working
- Inter-cell shared memory working
- ARM
 - Partly easier due to device tree description
 - Pitfall: shared resources (e.g. clock gate control on Banana Pi)
 - No publicly available reference setup so far
- Common limitation: no virtual consoles yet
 - Could be built on top of ivshmem
 - Or we add virtio compatibility to Jailhouse



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Inter-Cell Communication

ivshmem

- Shared r/w RAM region of two cells
- Signaling (MSIs)
- No messaging layer on top yet
- Why not... virtio, rpmsg, you-name-it...?
 - Must not share everything
 - Minimize copying
 - Minimize hypervisor effort
 - Avoid dynamic page remappings
- No perfect solution, need to find the least evil one
 - See also http://thread.gmane.org/gmane.linux.jailhouse/3479
 - BOF on virtio-based VM-to-VM communication at KVM Forum





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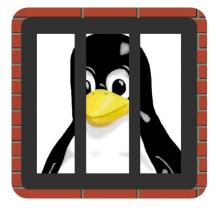
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Summary

Jailhouse provides clean AMP for Linux

- Full CPU isolation
- Minimal I/O latency
- Reduced to the minimum (goal: <10k LOC/arch)
- Jailhouse can run multiple Linux cells
 - Small Linux patch set for x86
 - Feasible with customizations on ARM
- Jailhouse is a community project
 - GPLv2, public development for 2 years
 - Significant contributions enabled / are enabling AMD64, ARMv7, and now ARMv8
 - You are invited to join!





Any Questions?

Thank you!

https://github.com/siemens/jailhouse

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